

Problem of the Month

Solution to Problem 4: January 2023

- (a) The positive integer 1 cannot be expressed as the sum of more than 1 positive integer, so if A compresses $[1 : 9]$, then A must contain the integer 1. We want A to be a list of positive integers of length four, and 1 is the smallest positive integer, so we can assume that $A = [1, a, b, c]$ where $a, b,$ and c are integers with $1 \leq a \leq b \leq c$.

The largest sum in $f(A)$ must be the sum of the three largest items in A , which is $a + b + c$, and since $f(A) = [1 : 9]$, we have $a + b + c = 9$.

Suppose that $a = 1$ and $b = 1$, then $a + b + c = 9$ implies that $c = 7$, but then $A = [1, 1, 1, 7]$, which does not satisfy $f(A) = [1 : 9]$ since, for example, 4 is not in $f(A)$. Therefore, a and b are not both 1.

Suppose that $a = 1$ and $b = 2$. Then $a + b + c = 9$ implies $c = 6$, so $A = [1, 1, 2, 6]$, but then $f(A)$ does not contain 5 so $f(A) \neq [1 : 9]$. Therefore, we cannot have $a = 1$ and $b = 2$.

Suppose that $a = 1$ and $b = 3$. Then $c = 5$, so $A = [1, 1, 3, 5]$. It can be verified that $f([1, 1, 3, 5]) = [1 : 9]$, so this gives one possible list A .

Suppose $a = 1$ and $b \geq 4$. Then $c \geq 4$ as well since $b \leq c$, but if $A = [1, 1, b, c]$ where $b \geq 4$ and $c \geq 4$, then $f(A)$ does not contain 3 which means $f(A) \neq [1 : 9]$.

So far, we have shown that $A = [1, 1, 3, 5]$ is the only list of the form we seek with $a = 1$ and $f(A) = [1 : 9]$.

We will now similarly examine the possibilities when $a = 2$.

If $a = 2$ and $b = 2$, then $c = 5$, so $A = [1, 2, 2, 5]$. It can be checked that $f(A) = [1 : 9]$ in this case.

If $a = 2$ and $b = 3$, then $c = 4$ and $A = [1, 2, 3, 4]$. It can be checked that $f(A) = [1 : 9]$ in this case.

If $a = 2$ and $b \geq 4$, then $a + b + c = 9$ implies that $c < 4$, but we are assuming that $b \leq c$, so this is impossible.

Therefore, the only possibilities with $a = 2$ are $A = [1, 2, 2, 5]$ and $A = [1, 2, 3, 4]$.

If $a \geq 3$, then $A = [1, a, b, c]$ has $b \geq 3$ and $c \geq 3$ as well, but this would prevent 2 from being in $f(A)$, so A cannot compress $[1 : 9]$ in this case.

Therefore, the only lists of length four that compress $[1 : 9]$ are $[1, 1, 3, 5]$, $[1, 2, 2, 5]$, and $[1, 2, 3, 4]$.

To see why no shorter list can compress $[1 : 9]$, we use the general observation that if A is a list of length k , then there are at most $2^k - 2$ integers in $f(A)$. This is because for each sum computed for $f(A)$, each of the k items in A is either included in the sum or it is not. This gives 2^k possible ways of computing a sum of some of the items in A . However, this count includes the sum of none of the items in A (all are excluded from the sum) and the sum of all of the items in A . Both of these sums are excluded from $f(A)$, so there are $2^k - 2$ ways to compute a sum to go in $f(A)$. We note that there could be multiple ways

to express the same integer in $f(A)$ as a sum of items in A , which is why we can only say there are *at most* $2^k - 2$ integers in $f(A)$ – in practice, there could be fewer than $2^k - 2$.

If $k \leq 3$, then $2^k - 2 \leq 2^3 - 2 = 6$, and so if A is a list of length at most three, then $f(A)$ has at most 6 integers. Therefore, $[1 : 9]$, which has nine integers, cannot be compressed by a list of length less than four.

- (b) At the end of the solution to part (a), we argued that if A is a list of length k , then there are at most $2^k - 2$ integers in $f(A)$. Since $2^6 - 2 = 62$ and $2^7 - 2 = 126$, we need $k \geq 7$ to achieve $2^k - 2 \geq 100$. Therefore, a list that compresses $[1 : 100]$ must have length at least seven. Now, we will provide a list of minimal length (seven) that compresses $[1 : 100]$.

Consider $A = [1, 2, 4, 8, 16, 32, 38]$. Note that 1 is in $f(A)$ and that the sum of all items in A is $1 + 2 + 4 + 8 + 16 + 32 + 38 = 101$. Since $f(A)$ has sums of some but not all of the items in A , the integers in $f(A)$ are at most 100. Rather than computing all of the sums, we will explain why $f(A) = [1 : 100]$ in a way that will give some insight for part (c).

We first consider the sums that are achievable without using the integer 38. The other integers in A are $1 = 2^0$, $2 = 2^1$, $4 = 2^2$, $8 = 2^3$, $16 = 2^4$, and $32 = 2^5$. The sums that can be obtained by adding some or all of these integers are exactly the integers from 1 through $63 = 2^6 - 1$. To get an idea of how this works, read about “binary expansions” or “binary representations” of integers. As an example, to represent 53 as a sum of powers of 2, first, find the largest power of 2 that is no larger than 53, which is 32. Then compute $53 - 32 = 21$ to get $53 = 32 + 21$. Now find the largest power of 2 that is no larger than 21, which is 16. Subtract to get $21 - 16 = 5$ or $21 = 16 + 5$. Now substitute to get $53 = 32 + 16 + 5$. Repeating this process, find the largest power of 2 that is no larger than 5, which is 4. Subtracting, $5 - 4 = 1$ so $5 = 4 + 1$, but what remains, 1, is a power of 2, so we get $53 = 32 + 16 + 4 + 1 = 2^5 + 2^4 + 2^2 + 2^0$.

The integers from 1 through 63 are all in $f(A)$. To write the integers from 64 through 100 as a sum of integers in A , notice that $100 - 38 = 62$, and so if $64 \leq m \leq 100$, then $m - 38 \leq 62$. To write such m as a sum of integers from A , compute $r = m - 38 \leq 62 < 63$, write r as a sum of the integers from A other than 38, then include 38 in the sum. For example, to see that 91 is in $f(A)$, compute $r = 91 - 38 = 53$, then use that $53 = 32 + 16 + 4 + 1$ to get that $91 = 1 + 4 + 16 + 32 + 38$.

The only thing left to check is that none of the sums described above require using all seven items in A . To see why this is not a concern, recall that the sum of all items in A is 101, so if we express an integer that is no larger than 100 as a sum of items from A , then it cannot possibly use every item in A .

- (c) The result of part (b) generalizes as follows. For every positive integer n , the minimum length of a list A that compresses $[1 : n]$ is $\lceil \log_2(n + 2) \rceil$. Notice that when $n = 100$, since $100 + 2 = 102$ is strictly between $2^6 = 64$ and $2^7 = 128$, we have $\lceil \log_2(100 + 2) \rceil = 7$, which agrees with the result from part (b).

We will prove that $\lceil \log_2(n + 2) \rceil$ is the minimum length of a list that compresses $[1 : n]$ and, in the process, prove that $[1 : n]$ is always compressible.

To see that $\lceil \log_2(n + 2) \rceil$ is the minimum length of a list that compresses $[1 : n]$ in general, we will show that a list of length less than $\lceil \log_2(n + 2) \rceil$ cannot possibly compress $[1 : n]$, and then we will construct a list of length exactly $\lceil \log_2(n + 2) \rceil$ that does compress $[1 : n]$.

Suppose A has length $k < \lceil \log_2(n+2) \rceil$. Since k and $\lceil \log_2(n+2) \rceil$ are both integers, this implies $k \leq \lceil \log_2(n+2) \rceil - 1$. For every integer x , it is true that $\lceil x \rceil - 1 < x \leq \lceil x \rceil$, so we conclude that $k \leq \lceil \log_2(n+2) \rceil - 1 < \log_2(n+2)$.

From $k < \log_2(n+2)$, we get $2^k < n+2$ or $2^k - 2 < n$. As argued earlier, a list A of length k has the property that there are at most $2^k - 2$ distinct integers in $f(A)$. Since $2^k - 2 < n$, we cannot have $f(A) = [1 : n]$ when A is a list of length $k < \lceil \log_2(n+2) \rceil$ since $[1 : n]$ contains n integers.

We have shown that if A compresses $[1 : n]$, then it must have length at least $\lceil \log_2(n+2) \rceil$. We will now produce a list of length $\lceil \log_2(n+2) \rceil$ that compresses $[1 : n]$. This requires explaining how to produce the list, then showing that the list has the correct length.

Suppose n is a positive integer. Define k to be the largest non-negative integer with the property that $2^k \leq n+1$ and define $m = n+2 - 2^k$. The list A consisting of the powers of 2 from 1 through 2^{k-1} together with m will compress $[1 : n]$ and have length exactly $\lceil \log_2(n+2) \rceil$. Notice that it is possible for m to be equal to one of the powers of 2 from 1 through 2^{k-1} . In this situation, the list A will include two copies of that power of 2.

Before verifying that the list described above does what is required, we will work through a couple of examples.

- When $n = 1$, we observe that $2^0 = 1$ and $2^1 = 2$ are the powers of 2 that are no larger than $n+1 = 2$, and so $k = 1$. Thus, $2^{k-1} = 1$ and $m = n+2 - 2^k = 1+2 - 2 = 1$, so the list is $A = [1, 1]$. Indeed $f([1, 1]) = [1]$.
- When $n = 100$, we get $k = 6$, so $2^{k-1} = 32$ and $m = n+2 - 2^k = 100+2 - 64 = 38$, so the list is $A = [1, 2, 4, 8, 16, 32, 38]$, which is exactly the list from part (b).

We will now show that list A has length $\lceil \log_2(n+2) \rceil$. The integer k is the largest non-negative integer with the property that $2^k \leq n+1$, and A contains $2^0, 2^1$, and so on up to 2^{k-1} , along with the integer m . This gives a total of $k+1$ items in A . Therefore, it suffices to show that with k chosen as described above, we have $\lceil \log_2(n+2) \rceil = k+1$.

The function \log_2 is increasing, meaning that if x and y are positive real numbers with $x < y$, then $\log_2(x) < \log_2(y)$. Using this fact along with the fact that $2^k \leq n+1$, we get that $k \leq \log_2(n+1) < \log_2(n+2)$. As well, k is the largest non-negative integer with the property that $2^k \leq n+1$, which means $n+1 < 2^{k+1}$.

Suppose $k+1 < \log_2(n+2)$. Then $2^{k+1} < n+2$. From above, we also have $n+1 < 2^{k+1}$, so we conclude that $n+1 < 2^{k+1} < n+2$. The quantities $n+1$ and $n+2$ are consecutive integers, so the integer 2^{k+1} cannot lie strictly between them. Therefore, it is impossible for $k+1 < \log_2(n+2)$, which means we must have $\log_2(n+2) \leq k+1$. Combining this inequality with $k < \log_2(n+2)$, we have $k < \log_2(n+2) \leq k+1$, and so we conclude that $\lceil \log_2(n+2) \rceil = k+1$.

It remains to show that A compresses $[1 : n]$. As discussed earlier, since list A contains the items $2^0, 2^1$, and so on up to 2^{k-1} , as well as at least one other item, m , $f(A)$ contains all of the integers from 1 through $2^{k-1+1} - 1 = 2^k - 1$. This is because these integers can be expressed using the sum of some or all of the powers of 2 from 1 through 2^{k-1} . These are exactly the integers that can be expressed without using m .

If we do use m , then we can express every integer from m through $m + 2^k - 1$ as a sum of

items in A . Since $m + 2^k - 1$ is the sum of all items in A , it is excluded from $f(A)$ and so we have that $f(A)$ contains all of the integers from m to $m + 2^k - 2$, and nothing larger. By definition, $m = n + 2 - 2^k$, so $m + 2^k - 2 = (n + 2 - 2^k) + 2^k - 2 = n$.

We have shown that $f(A)$ is the list consisting of the integers that are in $[1 : 2^k - 1]$ or $[m : n]$. To see that $f(A)$ is exactly $[1 : n]$, we need to show that $m \leq 2^k$. There might be overlap corresponding to multiple ways to express some integers in $[1 : n]$ in $f(A)$, but this is allowed.

Suppose $m > 2^k$. Since these quantities are integers, we must have $m \geq 2^k + 1$. By definition, $m = n + 2 - 2^k$, and so we get that $n + 2 - 2^k \geq 2^k + 1$, which can be rearranged to get $n + 1 \geq 2^k + 2^k = 2^{k+1}$. Therefore, we have $2^{k+1} \leq n + 1$, but k was chosen to be the largest integer with $2^k \leq n + 1$, so it is not possible that $2^{k+1} \leq n + 1$. Therefore, our assumption that $m > 2^k$ must be wrong, so we conclude that $m \leq 2^k$, as desired.

- (d) Fix a positive integer $k \geq 3$. We will show that for every integer $m \geq k$, $[m : m + k - 1]$ is not compressible.

To see this, suppose A is a list that compresses $[m : m + k - 1]$. Since $k \geq 3$, there are at least three items in $[m : m + k - 1]$. If A has only two items, then $f(A)$ has at most two integers since the allowable sums are just the two “singleton sums”. Therefore, A also contains at least three items. Note that since m is the smallest integer in $f(A)$, m must also be the smallest integer in A . (If r is in A , then the “singleton sum” r must also be in $f(A)$, so all integers in A must be at least m . Also, since there is nothing smaller than m in A , the only way to produce m in $f(A)$ is by the “singleton sum” m .)

Since $k \geq 3$, $m + 1$ is also in $f(A)$. If $m + 1$ is the sum of at least two items in A , then they must all be smaller than $m + 1$, but the only integer in A that is smaller than $m + 1$ is m . This would mean 1 must be in A , but this is impossible since $1 < k \leq m$ and m is the smallest element in A . Therefore, we must also have $m + 1$ in A , and so both m and $m + 1$ are in A .

Now recall that A has at least three items, so there is at least one element other than m and $m + 1$, and so $m + m + 1 = 2m + 1$ is in $f(A)$. Since $f(A) = [m : m + k - 1]$, we must then have $2m + 1 \leq m + k - 1$, which can be rearranged to get $m \leq k - 2$, which is impossible since $m \geq k$.

Therefore, it is not possible to compress $[m : m + k - 1]$ when $m \geq k$. This means that there are only finitely many m for which $[m : m + k - 1]$ is compressible since all such m must be at most k .

- (e) As mentioned in the hint, the answer is 39. This means $[5 : 39]$ is not compressible, but $[5 : k]$ is compressible for all $k \geq 40$. We will include a sketch of the proof here.

One can check that the lists $A = [5, 6, 7, 8, 9, 10]$, $B = [5, 5, 6, 6, 7, 8, 9]$, $C = [5, 5, 6, 7, 7, 8, 9]$, $D = [5, 5, 6, 7, 8, 8, 9]$, and $E = [5, 5, 6, 7, 8, 9, 9]$ compress the lists $[5 : 40]$, $[5 : 41]$, $[5 : 42]$, $[5 : 43]$, and $[5 : 44]$, respectively.

Now suppose a list A compresses $[5 : k]$ and suppose B is the list A with a 5 added to it. It can be shown that B compresses the list $[5 : k + 5]$. Since $[5 : 40]$ is compressible, this shows that $[5 : 45]$ is compressible. Since $[5 : 41]$ is compressible, $[5 : 46]$ is compressible. Since we have five consecutive values of k for which $[5 : k]$ is compressible ($k = 40$ through

$k = 44$), this reasoning can be used to show that $[5 : k]$ is compressible for all $k \geq 40$. Note that the lists to compress $[5 : k]$ given by this inductive process will not, in general, be as short as possible.

Now suppose a list A compresses $[5 : 39]$. It can be argued using reasoning similar to that from earlier parts that 5 must be in A , 5 is the smallest integer in A , and that the integers 6, 7, 8, and 9 also appear in A . Since the smallest integer in A is 5, the largest sum in $f(A)$ is 5 less than the sum of all items in A . Therefore, the sum of all items in A must be $39 + 5 = 44$. We already have 5, 6, 7, 8, and 9 in A which have a sum of $5 + 6 + 7 + 8 + 9 = 35$, and so the remaining items in A have a sum of $44 - 35 = 9$.

Next, note that using only the five items 5, 6, 7, 8, and 9, a sum of 10 is impossible. The list A cannot contain the integer 10 itself since we already determined that the remaining items in A have a sum of 9. It also cannot include the integers 1, 2, 3, or 4 since 5 is the smallest integer in A . Therefore, the 10 in $f(A)$ must come from the sum of two 5s, which means A must include a second 5.

We now have that A contains (at least) the items 5, 5, 6, 7, 8, and 9, which have a total of 40. Since the sum of all items in A is 44, there must be an additional item in A that is no larger than $44 - 40 = 4$. This is impossible, so we conclude that $[5 : 39]$ is not compressible.